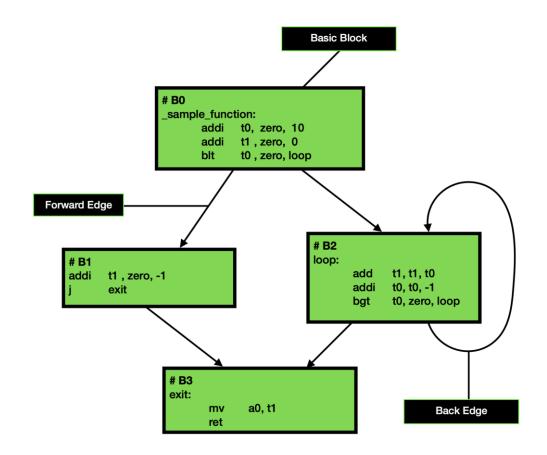
### **CMPUT 229**

Lab #5: Control Flow Graph

## What is a Control Flow Graph?

A Control Flow Graph (CFG) shows all the possible execution paths of a program.



### **Basic Blocks**

The nodes of a CFG are basic blocks.

A basic block is a segment of code where instructions execute in order from start to finish.

Once execution enters a block, all instructions in the block must be executed.

Only the first instruction of a basic block can be the target for a branch or jump instruction.

```
sample_function:
addi t0, zero, 10
addi t1, zero, 0
blt t0, zero, loop
```

```
addi t1, zero, -1
j exit
```

```
loop:
add t1, t1, t0
addi t0, t0, -1
bgt t0, zero, loop
```

```
exit:
mv a0, t1
ret
```

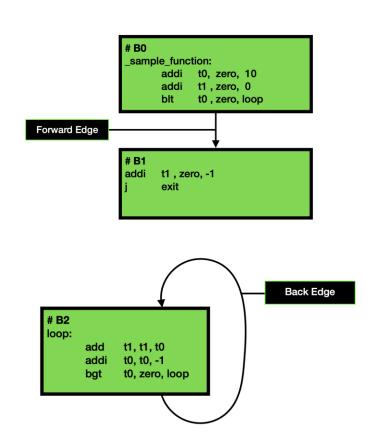
# Edges

Edges denote directed control flow between basic blocks.

There are two types of edges:

- forward edge
- back edge.

Branches and jumps redirect control flow, creating multiple execution paths within a program.

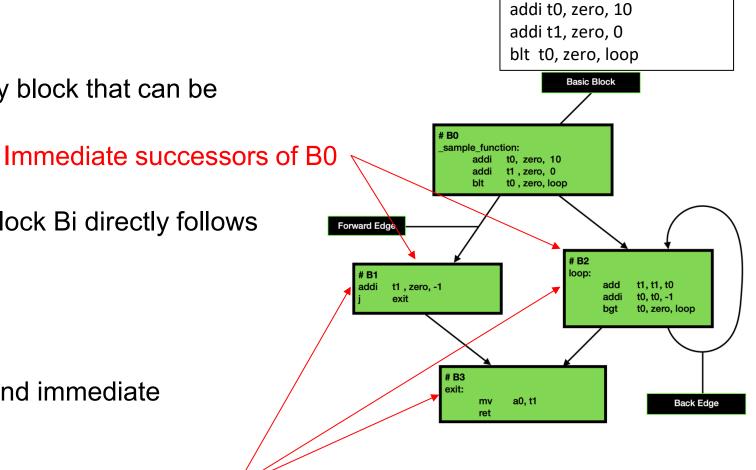


## Successors & Predecessors

A successor of a block Bi is any block that can be reached from Bi.

An immediate successor of a block Bi directly follows Bi.

The concepts of predecessor and immediate predecessor are analogous.



Successors of B0

#### cessors & Predecessors t0, zero, 10 t1, zero, 0 t0, zero, loop k Bi is any block that can be addi t0, zero, 10 # B2 addi t1, zero, 0 t1, zero, -1 t1, t1, t0 blt t0, zero, loop t0, t0, -1 t0, zero, loop Immediate successors of B0 # B3 sor of a block Bi directly follows exit: addi t1, zero, -1 Back Edge J exit The concepts of predecessor and immediate predecessor are analogous. Successors of B0

Lab #5: Control Flow Graph

## Assignment

### Overview

Goal: to generate control flow information for a single function.

These functions will populate data structures that represent control flow constructs:

- basic blocks, edges, successors, and predecessors.

Each function's implementation should accurately populate these custom data structures.

The **common.s** file initializes the data structures and runs tests on your implementation. It also helps to parse the populated data structures.

### Overview

- Goal: generate control flow information for a single function.
- Task: write functions that populate data structures that represent control flow constructs:
  - Basic blocks
  - Edges
  - Successors
  - Predecessors
- Resources: the common.s file:
  - initializes the data structures
  - runs tests on your implementation
  - helps to parse the populated data structures.

## A Sample Input Function

The **\_basicblocks** function below serves as an example of an intended input for the lab. It will be used as a reference to explain the custom data structures.

```
basicblocks:
             t0, zero
   mv
             t1, zero
   mv
   blez
             a0, done
  loop:
                t2, t0, 1
        andi
                t2, label1
        bnez
                t1, t1, t0
        add
                label2
        label1:
                     t1, t1, 1
             addi
        label2:
             addi
                     t0, t0, 1
                     a0, t0, loop
             bgt
   done:
       ret
```

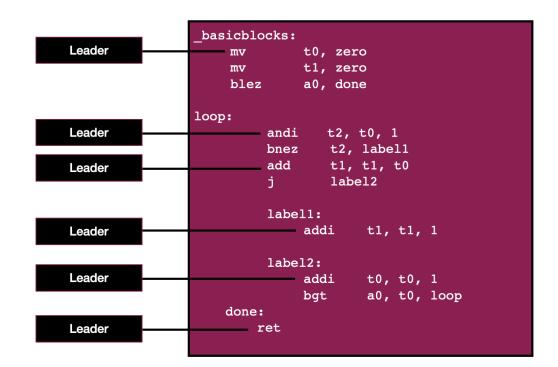
### **Basic Block Leaders**

Basic-block leaders:

The first statement in a function

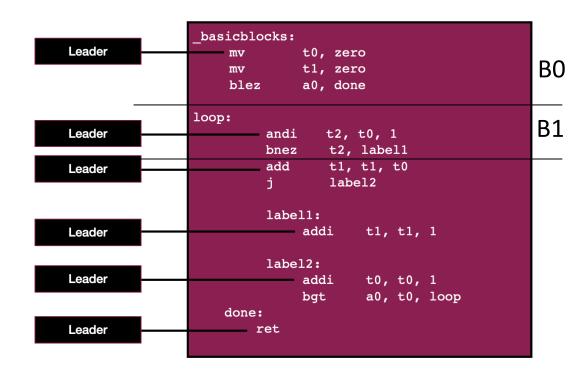
Any statement that is the target of a branch/jump

Any statement that immediately follows a branch or jump



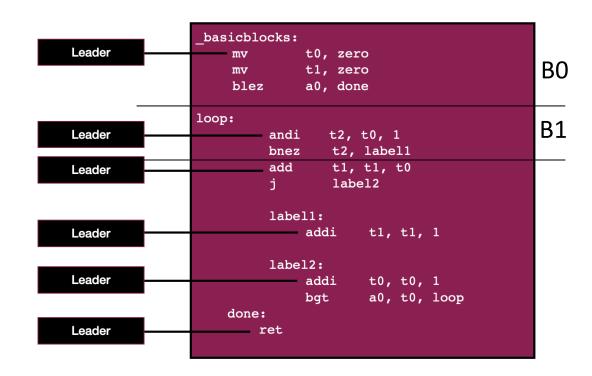
## Forming Basic Blocks

A basic-block is a leader and all instructions after the leader up to, but not including the next leader.



## Building the Control Flow Graph

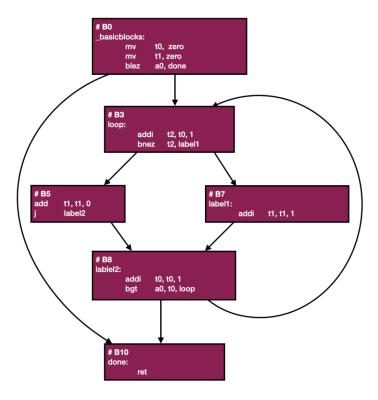
A basic-block is a leader and all instructions after the leader up to, but not including the next leader.



### **Basic Block Identifiers**

To uniquely identify blocks, we use the offsets (in words) of their leaders from the function's base address in memory.

Address	Hex	In	structions
00010000:	(000002b3)	add	t0, zero, zero
00010004:	(00000333)	add	t1, zero, zero
00010008:	(02a05063)	bge	zero, a0, 32
	(0012f393)		t2, t0, 1
00010010:	(00039663)	bne	t2, zero, 12
	(00530333)	add	t1, t1, t0
00010018:	(0080006f)	jal	zero, 8
0001001c:	(00130313)	addi	t1, t1, 1
00010020:	(00128293)	addi	t0, t0, 1
00010024:	(fea2c4e3)	blt	t0, a0, -24
00010028:	(00008067)	ret	



E.g., The second block starts 3 words away from the function's base address so its ID is 3.

## Pseudo Instructions and Labels

Pseudo-instructions in RISC-V simplify assembly code by representing higher-level operations.

		Edit Execute				
		Text Segment				
nents:						
Address	Code	Basic	Source			
	0×00400000	0x000002b3 add x5,x0,x0	4:	mv	t0, zero	
	0×00400004	0x00000333 add x6,x0,x0	5:	mv	t1, zero	
	0×00400008	0x02a05063 bge x0,x10,0x00000020	6:	blez	a0, done	
	0×0040000c	0x0012f393 andi x7,x5,1	9:	andi	t2, t0, 1	
	0×00400010	0x00039663 bne x7,x0,0x0000000c	10:	bnez	t2, label1	
	0×00400014	0x00530333 add x6,x6,x5	11:	add	t1, t1, t0	
	0×00400018	0x0080006f jal x0,0x00000008	12:	j	label2	
	0.0040001	0000120212 Add VE VE 1	15.	~44 · +1	+1 1	

For example, 'mv t0, zero' translates to 'add t0, zero, zero' during assembly, which copies the value 0 into the register t0 as well.

Labels are symbolic pointers to specific instructions used for branch or jump offsets and do not appear in the executable code. They are also used to simplify assembly.

# **Decoding Instructions**

RISC-V instruction formats are differentiated by their **opcode**, while **funct3** distinguishes instructions within each format.

31 30 25	24 21 20	19 15	14 12	11 8 7	6 0	
funct7	rs2	rs1	funct3	$\operatorname{rd}$	opcode	R-type
	1					-
imm[1]	1:0]	rs1	funct3	rd	opcode	I-type
[11 8]	0	1	C + O	. [4.0]	1	G .
imm[11:5]	rs2	rs1	funct3	imm[4:0]	opcode	S-type
:[10] :[10.5]	0	1	f 4 2	:[4.1] :[11]	1-	D 4
$[imm[12] \mid imm[10:5]$	rs2	rs1	funct3	imm[4:1]   imm[11]	opcode	B-type
	imm[31:12]			$\operatorname{rd}$	opcode	U-type
	111111[01:12]			Tu	opcode	C type
[imm[20]] $[imm[10]$	0:1] imm[11]	imm[1	9:12]	$\operatorname{rd}$	opcode	J-type

Decoding the fourth instruction (0x0012f393) of the **\_basicblocks** function as an example:

 $0 \times 0012f393 = 0000\ 0000\ 0001\ 0010\ 1111\ 0011\ 1001\ 0011$ 

opcode = 001 0011 (I-type instruction)

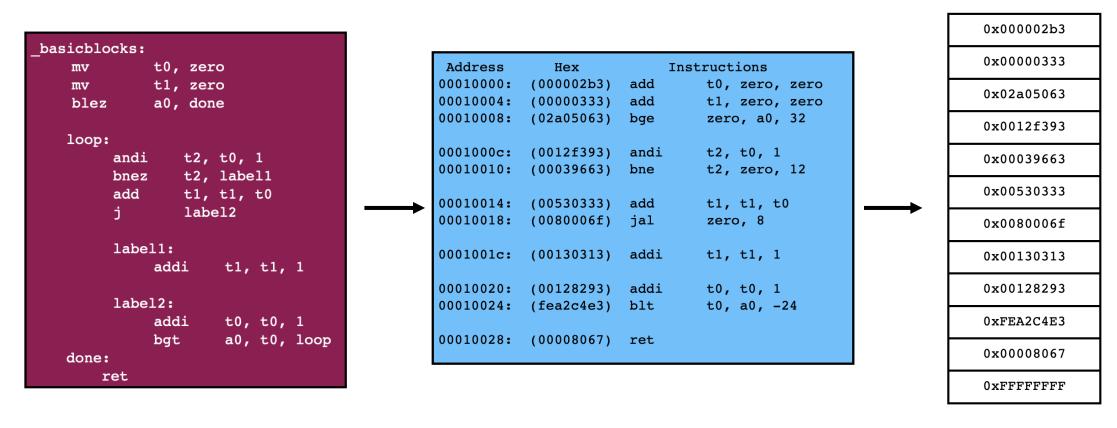
funct3 = 111 (andi instruction)

Lab #5: Control Flow Graph

### **Data Structures**

## instructionsArray

An array of words that contains the instructions of the input function.

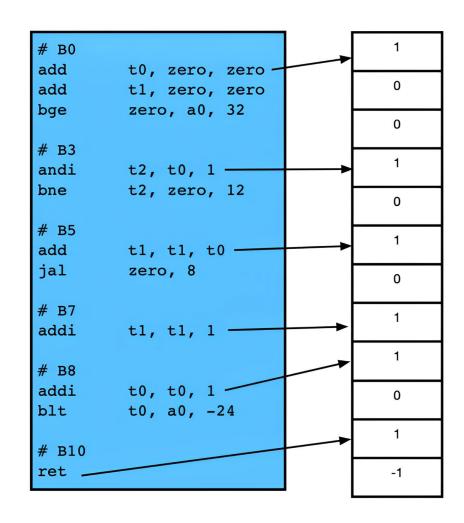


## leadersArray

An array of bytes representing the leadership status of each instruction.

Each byte corresponds directly to an instruction in the **instructionsArray**.

If an instruction is a leader, its corresponding byte will be 1; otherwise, it will be 0.

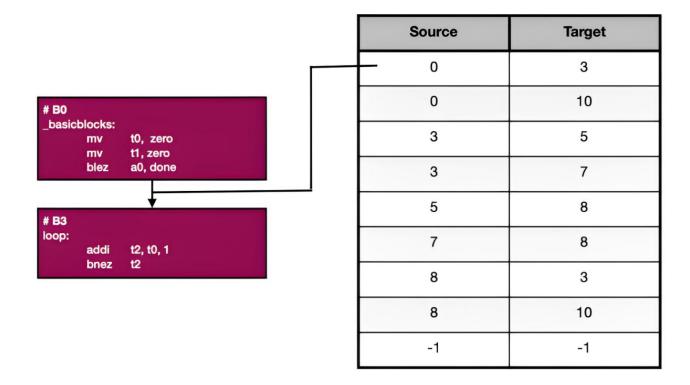


Ends with a sentinel value of -1.

# edgesList

A 2D byte array which stores all the edges within a function's CFG.

In an entry: the first byte indicates the source, and the second byte indicates the target.

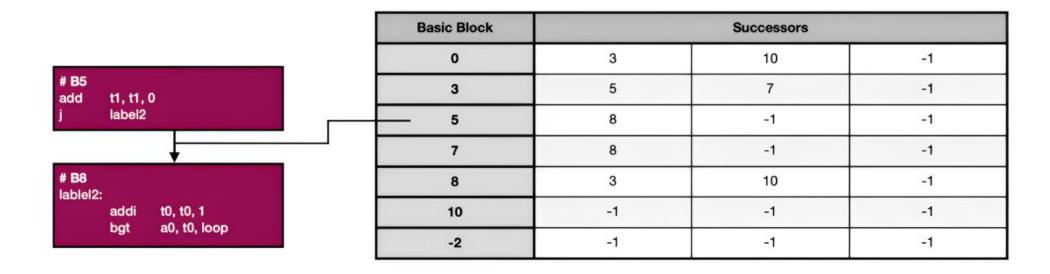


Ends with a sentinel value of -1 in both bytes.

### successorsTable

A 2D byte array storing immediate successors for each basic block in a function's CFG.

In an entry: the first byte indicates the block, followed by 2 bytes for its **successorsList**.



Each **successorsList** ends with -1, while the basic blocks column ends with -2 (indicating the table's end).

## predecessorsTable

A 2D byte array storing immediate predecessors for each basic block in a function's CFG.

In an entry: the first byte indicates the block, followed by 4 bytes for its **predecessorsList**.

		_ [	Basic Block			Predecessors				
# B3			0	-1	-1	-1	-1	-1		
loop: addi t2, t0, 1		3	0	8	-1	-1	-1			
	bnez t2		5	3	-1	-1	-1	-1		
	<b>+</b>		7	3	-1	-1	-1	-1		
# B5			8	5	7	-1	-1	-1		
add t1, t1, 0 j label2			10	0	8	-1	-1	-1		
			-2	-1	-1	-1	-1	-1		

Each **predecessorsList** ends with -1, while the basic blocks column ends with -2 (indicating the table's end).

Lab #5: Control Flow Graph

## **Function Signatures**

## getControlFlowGraph

#### **Description:**

The main function called from **common.s** to retrieve information about a function's control flow graph.

#### **Arguments:**

a0: a pointer to the instructionsArray

a1: a pointer to the leadersArray

a2: a pointer to the edgesList

a3: a pointer to the **successorsTable** 

a4: a pointer to the **predecessorsTable** 

#### **Returns:**

# getLeaders

### **Description:**

Identifies the basic block leaders within a function's instructions.

#### **Arguments:**

a0: a pointer to the instructionsArray

a1: a pointer to the **leadersArray** 

#### **Returns:**

1	0	0	1	0	1	0	1	1	0	1	-1

# getEdges

#### **Description:**

Creates a list of edges between basic blocks within a function's control flow graph.

#### **Arguments:**

a0: a pointer to the **instructionsArray** 

a1: a pointer to the leadersArray

a2: a pointer to the edgesList

#### **Returns:**

Source	Target
0	3
0	10
3	5
3	7
5	8
7	8
8	3
8	10
-1	-1

## getSuccessors

#### **Description:**

Creates a list of immediate successors for each basic block within a function's control flow graph.

#### **Arguments:**

a0: a pointer to the instructionsArray

a1: a pointer to the leadersArray

a2: a pointer to the **successorsTable** 

#### **Returns:**

Basic Block	Successors						
0	3	10	-1				
3	5	7	-1				
5	8	-1	-1				
7	8	-1	-1				
8	3	10	-1				
10	-1	-1	-1				
-2	-1	-1	-1				

## getPredecessors

#### **Description:**

Creates a list of immediate predecessors for each basic block within a function's control flow graph.

#### **Arguments:**

a0: a pointer to the instructionsArray

a1: a pointer to the leadersArray

a2: a pointer to the **predecessorsTable** 

#### **Returns:**

Basic Block	Predecessors							
0	-1	-1	-1	-1	-1			
3	0	8	-1	-1	-1			
5	3	-1	-1	-1	-1			
7	3	-1	-1	-1	-1			
8	5	7	-1	-1	-1			
10	0	8	-1	-1	-1			
-2	-1	-1	-1	-1	-1			

# getBranchImm (helper)

#### **Description:**

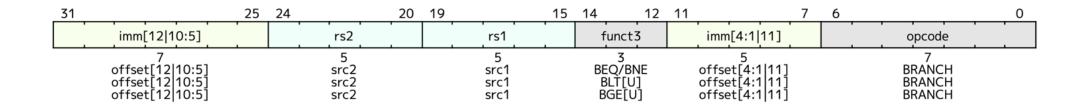
Takes a branch instruction (SB-type) and extracts the sign-extended immediate

#### **Arguments:**

a0: a **branch** instruction word

#### **Returns:**

a0: a sign extended immediate



# getJallmm (optional helper)

#### **Description:**

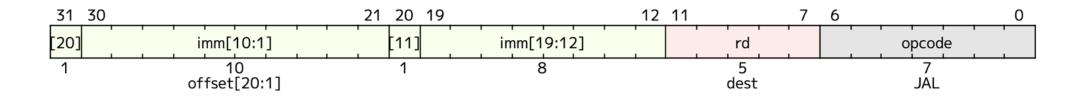
Takes a jump instruction (**UJ-type**) and extracts the sign-extended immediate

#### **Arguments:**

a0: a jal instruction word

#### **Returns:**

a0: a sign extended immediate



Lab #5: Control Flow Graph

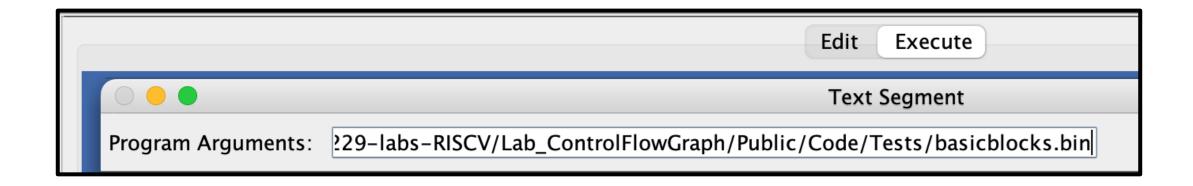
## Testing

## **Program Arguments**

We have provided some test inputs and expected outputs in the **Tests** folder.

One argument required for **controlFlowGraph.s**: the full path to the binary file of an assembled RISC-V function.

Ensure there are no quotation marks or spaces in the path.



### **Unit Tests**

The **common.s** file will run unit tests on the functions in **controlFlowGraph.s**.

Tests are hardcoded and do not use the file set as the program argument.

Check out the .data section in the common.s file to see how they are set up.

You can view the results in the "Run I/O" panel of RARS.

```
Messages Run I/O

-- Running tests for functions --

1: getLeaders -- [] Almost there!

2: getEdges -- [] Almost there!

3: getSuccessors -- [X] Great job!

4: getPredecessors -- [X] Great job!

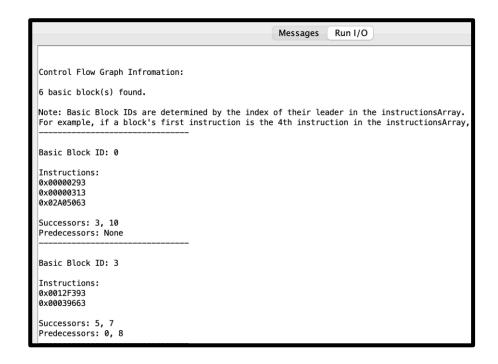
5: getControlFlowGraph --
```

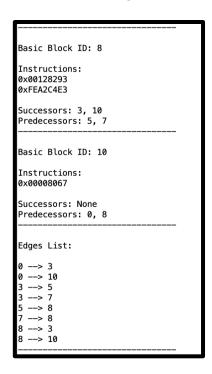
## Parsed Control Flow Graph

The **common.s** file also parses the data structures populated by **controlFlowGraph.s**.

All structures must end with a sentinel value to be parsed properly.

You can view the results in the "Run I/O" panel of RARS following the unit tests.





### Tests Folder

There are three tests: basicblocks, nestedloop, singleinstruction.

Each test has the following:

- A .s file containing the assembly code of the input function.
- A .bin (binary) file containing the assembled function (program argument).
- A **.out** file containing the correct output for the test.

## **Creating Tests**

Unit tests and provided tests are not extensive.

Create your own tests to ensure your lab handles corner cases as well.

To create a test, do the following:

- Write the test function in assembly file (e.g., test.s).
- Create the binary file for your function i.e., execute the command "rars a dump .text Binary test.bin test.s".
- You can now execute **controlFlowGraph.s** with the program argument set to the full path of **test.bin**.

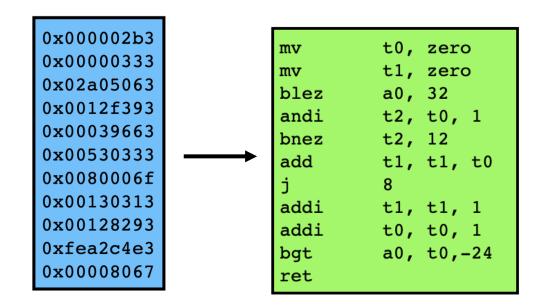
Lab #5: Control Flow Graph

## Disassembler

### What is a Disassembler?

RARS is a RISC-V Assembler that translates RISC-V instructions to executable binary.

A Disassembler does the opposite.



For this lab, we will use an Open-Source RISC-V Disassembler (<a href="https://github.com/michaeljclark/riscv-disassembler">https://github.com/michaeljclark/riscv-disassembler</a>)

### How to Use the Disassembler

There is a **Disassembly** folder in the **Code** folder.

There is a file called "print-instructions.c" in this folder that prints the equivalent instructions for hexadecimal words.

First, compile "print-instructions.c" i.e., execute the command "gcc print-instructions.c"

Next, create a text file containing hexadecimal instructions. The file should look like this:

## How to Use the Disassembler (cont.)

Once you have an executable for print-instructions.c (a.out) and a text file with hexadecimal instructions (example.txt), execute "./a.out example.txt"

Here is the disassembled instructions from **example.txt**:

```
./a.out example.txt
0000000000010000:
                   00100413
                                      addi
                                                     s0,zero,1
0000000000010004:
                   02850433
                                      mul
                                                     s0,a0,s0
                                      addi
000000000010008:
                   fff50513
                                                     a0,a0,-1
000000000001000c:
                   00050463
                                                     a0,8
                                      beqz
                                                                                        0x10014
                                                     -12
0000000000010010: ff5ff06f
                                                                                      # 0x10004
                                                     a0, zero, s0
0000000000010014:
                   00800533
                                      add
0000000000010018:
                   00008067
                                      ret
```

### Notes on the Disassembler

The disassembler translates **addi t0, t0, 0** to **mv t0, zero**. Keep this in mind to avoid confusion with **addi** instructions.

The disassembler will not translate a sentinel value (0xFFFFFFF) as expected.

### What to Submit?

A single file, called **controlFlowGraph.s**.

**Keep** the file in the **Code** folder of the git repository.

Do not modify the name of any function.

Do not remove the CMPUT 229 Student Submission License.

Do not modify the line .include "common.s".

Do not modify the common.s file.

**Push** your repository to GitHub before the deadline.

Lab #5: Control Flow Graph

## Good Luck!